

Bargaining Foundations of the Median Voter Theorem

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Outline of talk

- brief introduction
- bargaining literature
- model
- voting subgames
- main result
- intuition for proof
- extensions
- quick conclusion

The bargaining literature

Parameters	Rubinstein	Binmore
players	$n = 2$	$n = 2$
alternatives	Δ/\mathfrak{R}	Δ/\mathfrak{R}
voting rule	U/M	U/M
proposer	deterministic	random
status quo	bad	bad

Parameters	Baron-Ferejohn	Banks-Duggan
players	$n > 2$	$n > 2$
alternatives	Δ	\mathfrak{R}
voting rule	M	M
proposer	random	random
status quo	bad	whatever

The bargaining literature (cont.)

Issues	R	B	B-F	B-D
existence of stationary eq.	Y	Y	Y	Y
delay in stationary eq.	N	N	N	N*
uniqueness of stationary eq.	Y*	Y*	Y	N
asymptotic mvt	n/a	n/a	n/a	Y
folk thm	N*	N*	Y	

Elements of the model

N	set of n agents (n odd)*
X	set of alternatives (a compact interval)
$u_i: X \rightarrow \mathfrak{R}$	stage utility (cont., strictly concave)*
\tilde{x}_i	i 's ideal point (median is \tilde{x}_k)
q	status quo (in X)*
δ	discount factor (common, less than one)*
ρ_i	recognition prob. (positive, fixed)*

The bargaining protocol in period t

- an agent i is selected with probability ρ_i to propose, say, x
- each agent j votes accept (a) or reject (r)
- if $\#\{j \mid v_j = a\} > n/2$, then the game ends with outcome (x, t) , and payoffs are

$$(1 - \delta)^{t-1}u_j(q) + \delta^{t-1}u_j(x)$$

- otherwise, the game continues to period $t + 1$ and is repeated
- if the game continues ad infinitum, then payoffs are $u_j(q)$

The bargaining protocol (cont.)

- Think of payoffs from (x, t) as the discounted sum of the flow of payoffs

$$\underbrace{u_i(q) \ u_i(q) \ \cdots \ u_i(q)}_{t-1 \text{ periods}} \quad \underbrace{u_i(x) \ u_i(x) \ u_i(x) \ \cdots}_{\text{ad infinitum}}$$

Continuation lotteries

- We can summarize future play by a “continuation lottery.”
- For example, suppose in odd periods, there is a $1/3$ chance that x passes and $2/3$ chance of delay; in even periods, there is a $1/4$ chance of x' and $3/4$ chance of delay.

Continuation Lotteries (cont.)

- Then agent i 's expected payoff U_i at the beginning of period 1 can be written as

$$U_i = \alpha u_i(x) + \beta u_i(x') + \gamma u_i(q)$$

where

$$\alpha = \frac{2}{6 - 3\delta^2}$$
$$\beta = \frac{\delta}{6 - 3\delta^2}$$
$$\gamma = \frac{4 - \delta - 3\delta^2}{6 - 3\delta^2}$$

and

$$\alpha + \beta + \gamma = 1.$$

- That is, $U_i = E_\lambda[u_i]$, where λ is the continuation lottery on X .

Voting subgames

- In stationary equilibria, the voting game reduces to a binary vote, and stage-game weak dominance is effective.
- Suppose x has been proposed.

	a	r
a	x	x
r	x	λ
	a	

	a	r
a	x	λ
r	λ	λ
	r	

If row strictly prefers x to continuing, then voting r is weakly dominated in the stage game.

- In fact, x survives elimination if and only if it is weakly majority-preferred to continuing.

Voting subgames (cont.)

- Not so when stationarity is dropped.

	<i>a</i>	<i>r</i>
<i>a</i>	<i>x</i>	<i>x</i>
<i>r</i>	<i>x</i>	λ_1
	<i>a</i>	

	<i>a</i>	<i>r</i>
<i>a</i>	<i>x</i>	λ_2
<i>r</i>	λ_3	λ_4
	<i>a</i>	<i>r</i>

- Suppose payoffs are as follows.

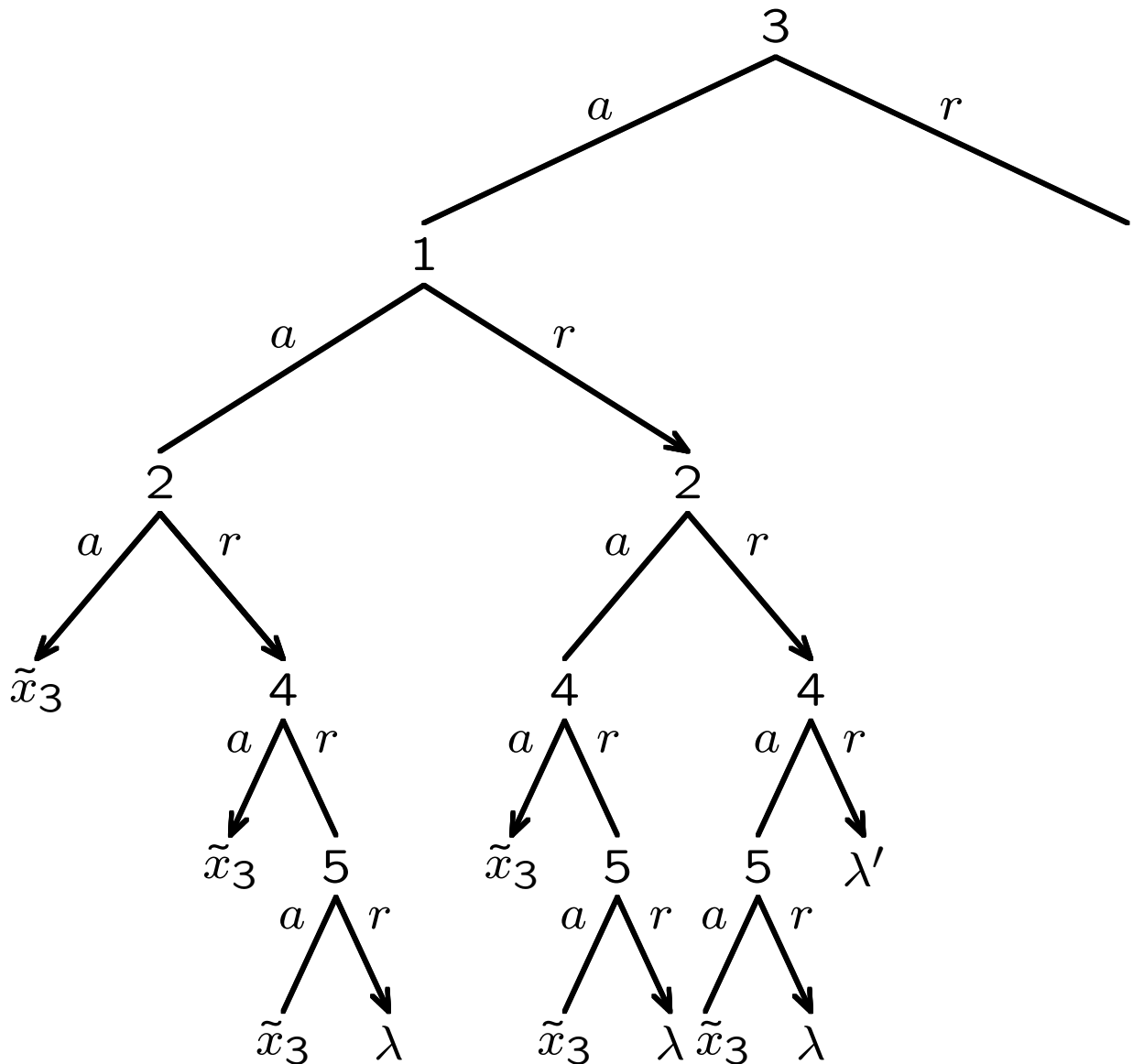
row	column	matrix
λ_1	λ_2	λ_3
<i>x</i>	<i>x</i>	<i>x</i>
λ_4	λ_4	λ_4
λ_2	λ_3	λ_1

Then λ_4 is a strict (therefore undominated) Nash equilibrium outcome, but all agents strictly prefer *x*.

Voting subgames (cont.)

- We model voting as sequential: this allows us to drop the stage-game weak dominance refinement and leaves stationary equilibrium outcomes unchanged.
- A key property we want: the median \tilde{x}_k passes if it is proposed.
- For this, we assume:
 - The order of voting is determined randomly before and/or after the proposer is determined.
 - With positive, fixed* probability, the median voter votes first, and voting then alternates from either side of the median.

- The problem: Let ideal points of 1, 2, 3, 4, 5 be in order of voter indices. Suppose the median \tilde{x}_3 is proposed and 3 has voted $a \dots$



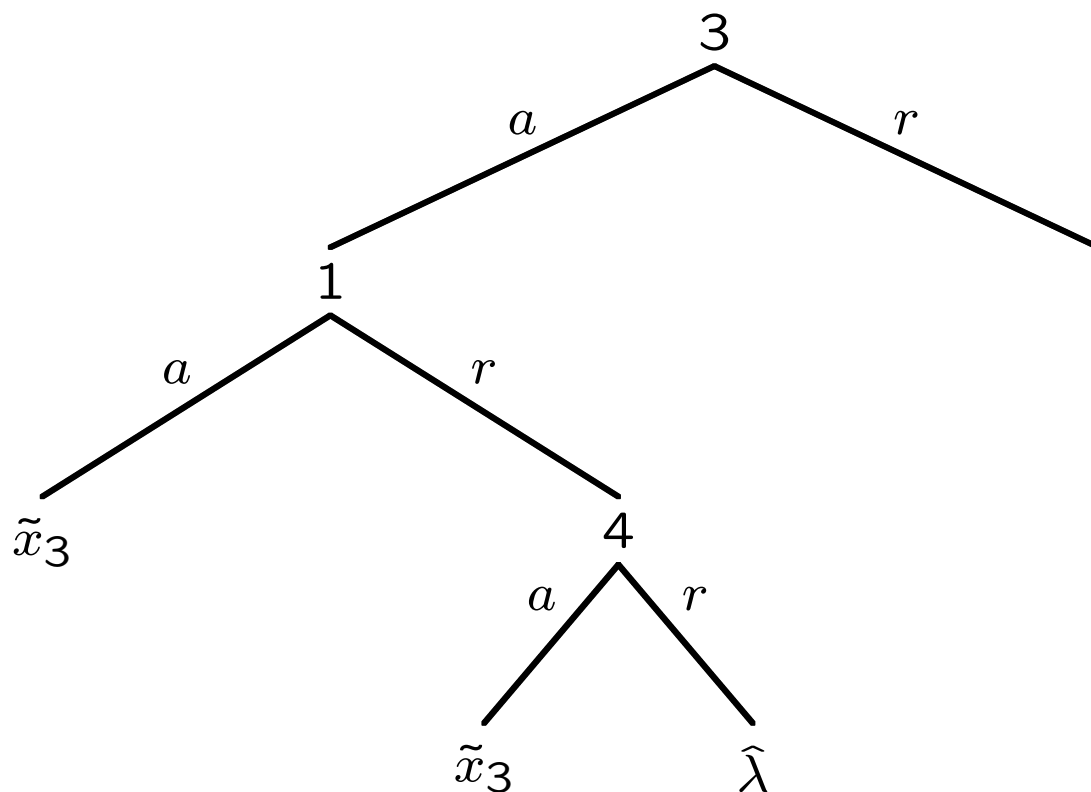
1	2	4	5
λ'	λ'	\tilde{x}_3	λ
\tilde{x}_3	\tilde{x}_3	λ'	\tilde{x}_3
		λ	

Voting subgames (cont.)

- By concavity, if agent 5 weakly prefers λ to the median \tilde{x}_3 , and if agent 2 has the same weak preference, then the mean of λ must be equal to \tilde{x}_3 .
- By strict concavity, the continuation lottery λ must have zero variance, so it is the point mass on \tilde{x}_3 .
- Therefore, the only possible equilibrium outcome of the first two 2:5 voting subgames is the median.

Voting subgames (cont.)

- Then the game reduces to...



... and the same logic applies.

- Therefore, the median voter can (and will) obtain her ideal point by voting a in equilibrium.

Subgame perfect equilibria

- A strategy σ is a pure strategy subgame perfect equilibrium if there do not exist an agent, a history, and a deviation from σ that increases that agent's expected payoff following that history.
- Let $X(\delta)$ consist of alternatives x such that:
 - there exists a proposer history h and a pure strategy subgame perfect equilibrium σ such that x passes with positive probability from h in σ .
- Let $V_i(\delta)$ consist of payoffs r such that:
 - there exists a proposer history h and a pure strategy subgame perfect equilibrium σ such that i 's expected payoff is r from h in σ .

The main result

- Given a sequence of sets $\{Y^m\}$ in Euclidean space and an element x , we write $Y^m \rightarrow x$ if $\sup_{y \in Y^m} \|y - x\| \rightarrow 0$.
- Theorem: Let $\{\delta^m\}$ be a sequence of discount factors converging to one. Then
 - $X(\delta^m) \rightarrow \tilde{x}_k$,
 - for all $i \in N$, $V_i(\delta^m) \rightarrow u_i(\tilde{x}_k)$.
- That is,
 - subgame perfect equilibrium outcomes converge to the median ideal point,
 - equilibrium delay becomes negligible.

Intuition for proof (cont.)

- Lemma: Suppose $x \in X(\delta^m)$ is proposed in some equilibrium and passes with positive probability. Then there is an agent i^m in C^m who weakly prefers x to the continuation lottery from rejecting the proposal.
- Apply the lemma with $x = \bar{x}^m$ to get an agent i^m from each C^m who weakly prefers \bar{x}^m to the continuation lottery from rejection. But...
- Lemma: In every subgame perfect equilibrium, if the median \tilde{x}_k is proposed, then it passes with positive probability (bounded above zero).

Intuition for proof (cont.)

- As a consequence, agent i^m can obtain the median with positive probability when selected to propose.
- So why would the agent i^m be willing to vote for his/her worst alternative rather than reject it and at least obtain the median \tilde{x}_k with positive probability?
- It must be that the status quo is inferior to \bar{x}^m , and that delay makes the agent worse off.
- But we show that, as the agent becomes patient, delay becomes inconsequential, a contradiction.

Extensions

- Even number of agents: okay, but we need one of the median voters to “break ties.”
- Concavity: strict concavity can be weakened to allow for piece-wise linear.
- Bad status quo: we can take the status quo out of X as long as every agent strictly prefers the median to the status quo.
- Heterogeneous discount factors: we just need discount factors to converge to one at a “uniform enough” rate, i.e.,

$$\frac{\ln(\delta_i^m)}{\ln(\delta_j^m)} \rightarrow 1$$

for all $i, j \in N$.

Extensions (cont.)

- Variable recognition probabilities: These can vary arbitrarily with histories, as long as each agent's probability is bounded strictly above zero.
- Variable probability of alternating voting order: Can vary arbitrarily with histories, as long as voting alternates with probability bounded strictly above zero.
- Mixed strategy equilibria: mixed proposal strategies are not a problem; mixed voting strategies introduce some complications, but these should not be a problem.

Quick conclusion

- In contrast to the distributive model of bargaining, an anti-folk theorem holds for the unidimensional model.
- The asymptotic median voter theorem for stationary equilibria extends quite generally.
- Delay becomes negligible as agents become patient.
- The median voter theorem for small groups is pretty “safe.”